

## Assignment 4 Solution

*Instructor: Sourav Chakraborty**Scribe: Ankan Kumar Das***2. Prove or Disprove the set of asymptotic relations.**

(a)  $2.9^{\log_2 n} = \Theta(n^{\log_2 3})$

$$\lim_{n \rightarrow \infty} \frac{n^{\log_2 3}}{2.9^{\log_2 n}} = \lim_{n \rightarrow \infty} \frac{n^{\log_2 3}}{n^{\log_2 2.9}} \sim \lim_{n \rightarrow \infty} n^{0.049} \rightarrow \infty$$

**Disproved**

(b)  $\log \log n = \Omega((\log \log \log n)^{\log \log \log n})$

Let  $x = \log \log n$  and  $y = \log \log \log n^{\log \log \log n}$ 

Hence,  $\lim_{n \rightarrow \infty} \frac{\log y}{\log x} = \lim_{n \rightarrow \infty} \log \log \log \log n \rightarrow \infty$

**Disproved**

(c)  $n^4 \sim (1 - \frac{1}{n})^n n^3$

Let  $x = n^4$  and  $y = (1 - \frac{1}{n})^n n^3$ Then  $\log x = 4 \log n$  and  $\log y = n \log(1 - 1/n) + 3 \log n$ 

$$\lim_{n \rightarrow \infty} \frac{\log y}{\log x} = \lim_{n \rightarrow \infty} \frac{n \log(1-1/n) + 3 \log n}{4 \log n} = \frac{3}{4}$$

Hence, they are not asymptotically equal.

(d)  $2^{(\log n) - (\log \log n)} \sim 2^{(1-1/n) \log n}$

Let  $x = 2^{(\log n) - (\log \log n)}$  and  $y = 2^{(1-1/n) \log n}$ Then  $\log x = (\log n - \log \log n) \log 2$  and  $\log y = (1 - 1/n) \log n \log 2$ 

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{\log x}{\log y} &= \lim_{n \rightarrow \infty} \frac{1/n - 1/(n \log n)}{1/n - 1/n^2 + \log n/n^2} \quad [\text{Using L'Hopital's Rule}] \\ &= \lim_{n \rightarrow \infty} \frac{1 - 1/\log n}{1 - 1/n + \log n/n} \\ &= 1 \end{aligned}$$

Hence, they are asymptotically equal.

**3. Prove that**  $\binom{2n}{n} \sim \frac{4^n}{\sqrt{\pi n}}$

$$\binom{2n}{n} = \frac{(2n)!}{n!n!} \sim \frac{\sqrt{2\pi(2n)}\left(\frac{2n}{e}\right)^{2n}}{(\sqrt{2\pi n}\left(\frac{n}{e}\right)^n)^2} \quad [\text{Using Stirling's Approximation}]$$

$$\text{Hence, } \binom{2n}{n} \sim \frac{4^n}{\sqrt{\pi n}}$$

**4. For the following pairs give correct asymptotic relations**

(a)  $(\log n)^a$  and  $n^b$

Claim:  $(\log n)^a \leq cn^b$

For the claim to be true we need to find  $c$  s.t.  $c \geq \frac{(\log n)^a}{n^b}$

$$\text{Let, } f(x) = \frac{(\log x)^a}{x^b}$$

$$\text{Therefore } f'(x) = \frac{a(\log x)^{a-1}}{x^{b+1}} - \frac{b(\log x)^a}{x^{b+1}}$$

$$\text{and } f''(x) = \frac{a(a-1)(\log x)^{a-2} - a(2b+1)(\log x)^{a-1} + b(b+1)(\log x)^a}{x^{b+2}}$$

$$f'(x) = 0 \Rightarrow \log x = \frac{a}{b}, \text{ i.e. } x = e^{\frac{a}{b}}$$

$$f''(e^{\frac{a}{b}}) \leq 0 \quad \forall b \in \mathbb{R}$$

Therefore  $f(x)$  is maximum at  $x = e^{\frac{a}{b}}$

If we choose  $c = f(e^{\frac{a}{b}})$  then it will satisfy our claim.

$$\text{Hence } (\log n)^a = O(n^b)$$

(b)  $2^{n \log_2 n}$  and  $10n!$

$$2^{n \log_2 n} = n^n \text{ and } n! \sim \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$$

$$\text{Now } \lim_{n \rightarrow \infty} \frac{n^n}{10n!} = \lim_{n \rightarrow \infty} \frac{e^n}{10\sqrt{2\pi n}} \rightarrow \infty$$

$$\text{Therefore, } 2^{n \log_2 n} = \omega(10n!)$$

(c)  $\sqrt{n}$  and  $(\log_2 n)^5$

$$\lim_{n \rightarrow \infty} \frac{\sqrt{n}}{(\log_2 n)^5} = \lim_{n \rightarrow \infty} \frac{\frac{1}{2\sqrt{n}}}{\frac{5(\log_2 n)^4}{n \log_e 2}} \quad [\text{Using L'Hopital's Rule}]$$

$$= \frac{\log_e 2}{10} \lim_{n \rightarrow \infty} \frac{\sqrt{n}}{(\log_2 n)^4} = \frac{(\log_e 2)^2}{80} \lim_{n \rightarrow \infty} \frac{\sqrt{n}}{(\log_2 n)^3} = \dots = \frac{(\log_e 2)^5}{2^{55}} \lim_{n \rightarrow \infty} \sqrt{n} \rightarrow \infty$$

$$\text{Therefore, } \sqrt{n} = \omega((\log_2 n)^5)$$

(d)  $n^2 / \log_2 n$  and  $(n \log_2 n)^4$

$$\lim_{n \rightarrow \infty} \frac{n^2 / \log_2 n}{(n \log_2 n)^4} = \lim_{n \rightarrow \infty} \frac{1}{n^2 (\log_2 n)^5} \rightarrow 0$$

Therefore,  $n^2 / \log_2 n = o(n \log_2 n)^4$

(e)  $\log_2 n$  and  $\log_2 66n$

$$\lim_{n \rightarrow \infty} \frac{\log_2 n}{\log_2 66n} = \lim_{n \rightarrow \infty} \frac{\log_2 n}{\log_2 n + \log_2 66} \rightarrow 1$$

Therefore,  $\log_2 n = \Theta(\log_2 66n)$

(g)  $n^2$  and  $n(\log_2 n)^{15}$

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{n^2}{n(\log_2 n)^{15}} &= \lim_{n \rightarrow \infty} \frac{n}{(\log_2 n)^{15}} = \frac{\log_e 2}{15} \lim_{n \rightarrow \infty} \frac{n}{(\log_2 n)^{14}} \quad [\text{Using L'Hopital's Rule}] \\ &= \dots = \frac{(\log_e 2)^{15}}{15!} \lim_{n \rightarrow \infty} n \rightarrow \infty \end{aligned}$$

Therefore,  $n^2 = \omega(n(\log_2 n)^{15})$

5. Find approximate value of  $\binom{n}{n/3}$

$$\begin{aligned} \binom{n}{n/3} &= \frac{n!}{(n/3)!(2n/3)!} \sim \frac{\sqrt{2\pi n}(n/e)^n}{\sqrt{2\pi(n/3)}(n/3e)^{n/3} \sqrt{2\pi(2n/3)}(2n/3e)^{2n/3}} \quad [\text{Using Stirling's Approximation}] \\ &= \frac{3^{n+1}}{2^{2n/3} \sqrt{4\pi n}} \end{aligned}$$

6. Solve the following recurrences

(a)  $T(n) = 2T(\lceil n/2 \rceil) + 5$

Claim:  $T(n) \leq cn - d$  for some +ve  $c$  and  $d$  which will be fixed later

**Proof:**

From the base case  $T(2) = 7$  condition we have  $2c - d \geq 7$

$$\begin{aligned} T(n) &= 2T(\lceil n/2 \rceil) + 5 \\ &\leq 2c\lceil n/2 \rceil - 2d + 5 \quad [\text{According to I.H.}] \\ &\leq 2cn/2 + c + 5 - 2d \\ &= cn - 2d + c + 5 \\ &\leq cn - d \quad [\text{For } c + 5 - d \geq 0] \end{aligned}$$

$c = 12$  and  $d = 17$  satisfies both the condition.

Hence,  $T(n) \leq 12n - 17 \forall n \geq 2$

$$(b) T(n) = T(n-1) + T(n-2) + 15$$

Claim:  $T(n) \leq c2^n \quad \forall n \geq n_0$  where  $n_0$  will be decided later

**Proof:**

We have  $T(2) = 17$  and  $T(3) = 33$ . Hence  $4c \geq 17$  and  $8c \geq 33$  i.e.  $c \geq 5$  satisfies both the condition.

$$\begin{aligned} T(n) &= T(n-1) + T(n-2) + 15 \\ &\leq c2^{n-1} + c2^{n-2} + 15 && \text{[According to I.H.]} \\ &= c2^n + 15 - c(1 + 1 + 2 + 2^2 + 2^3 + \dots + 2^{n-3}) \\ &= c2^n + 15 - c2^{n-2} \\ &\leq c2^n && \text{when } 15 - c2^{n-2} \leq 0 \end{aligned}$$

Therefore,  $c2^{n-2} \geq 15 \Rightarrow 2^{n-2} \geq 3 \Rightarrow n \geq 4$

For  $n < 4$  our claim is true as we can check manually and from principle of mathematical induction claim is true for  $n \geq 4$  as well

$$(c) T(n) = 2T(\lceil n/3 \rceil) + n^2$$

Claim:  $T(n) \leq cn^2$  for some +ve  $c$  which will be decided later

**Proof:**

We have  $T(1) = 1, T(2) = 5, T(3) = 11, T(4) = 26$ .  $c = 2$  satisfies our assumption for these values of  $n$ .

$$\begin{aligned} T(n) &= 2T(\lceil n/3 \rceil) + n^2 \\ &\leq 2c\lceil n/3 \rceil^2 + n^2 && \text{[According to I.H.]} \\ &\leq 2c(n/3 + 1)^2 + n^2 \\ &= 2c(n^2/9 + 2n/3 + 1) + n^2 \\ &= cn^2 - (7cn^2/9 - 4n/3 - 2c - n^2) \\ &\leq cn^2 && \text{when } 7cn^2/9 - 4n/3 - 2c - n^2 \geq 0 \end{aligned}$$

$$7cn^2/9 - 4n/3 - 2c - n^2 \geq 0$$

$$\Rightarrow 14n^2/9 - 4n/3 - 4 - n^2 \geq 0$$

$$\Rightarrow 5n^2 - 12n - 36 \geq 0$$

For  $n \geq 5$  the above condition satisfies.

Hence by Principle of Mathematical Induction  $T(n) \leq 2n^2 \quad \forall n \geq 1$