

Lecture 1: Mid sem solutions(1 - 6)

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1. Let $x \in \mathbb{R}$ and $x > -1$. Prove that $(1+x)^n \geq 1+nx$ for all natural numbers n

Sol: We proceed by induction on $n \in \mathbb{N}$.

Let $P(n)$ be the proposition that $(1+x)^n \geq 1+nx$ for all natural numbers n , whenever $x > -1$, $x \in \mathbb{R}$

Base case: $P(1)$ is true

Induction Hypothesis: Let $P(k)$ be true for $k \in \mathbb{N}$.

Inductive Step: For $n=k+1$,

$$\begin{aligned} (1+x)^{k+1} &= (1+x)^k(1+x) \\ &\geq (1+kx)(1+x) \text{ [By Induction Hypothesis and because } (1+x) \geq 0] \\ &\geq 1+(k+1)x+kx^2 \\ &\geq 1+(k+1)x \text{ as } x^2 \geq 0 \end{aligned}$$

Hence by induction, our proposition is true for all $n \in \mathbb{N}$.

2. Prove or disprove the following set of asymptotic relations:

- (a) $(2.9)^{\log_2 n} = \Theta(n^{\log_2 3})$
- (b) $\log \log n = \Omega((\log \log \log n)^{\log \log \log n})$
- (c) $n^4 \sim (1 - \frac{1}{n})^n$
- (d) $2^{\log n - \log \log n} \sim 2^{(1 - \frac{1}{n}) \log n}$
- (e) $n^{10(\log \log n)^{100}} = \Theta((\log n)!)$

Sol:

(a) We have $(2.9)^{\log_2 n} = n^{\log_2(2.9)}$. Since $\log_2(2.9) \leq \log_2 3$, so $\frac{(2.9)^{\log_2 n}}{n^{\log_2 3}} = n^{\log_2(2.9) - \log_2 3}$ which tend to 0 as n goes to infinity. So $(2.9)^{\log_2 n} = o(n^{\log_2 3})$ and hence the given statement is not true

(b) Let's take, $x = \log \log \log n$.

Now,

$$e^x = o(x^x) \quad \forall \text{ increasing function } x.$$

$$\therefore \log \log n = O((\log \log \log n)^{\log \log \log n})$$

, disproving the given statement.

(c) $\lim_{n \rightarrow \infty} (1 - \frac{1}{n})^n = e^{-1}$ while $\lim_{n \rightarrow \infty} n$ is divergent. Thus we disprove the given statement.

(d) Let $x = \log n - \log \log n$ and $y = (1 - \frac{1}{n}) \log n$.

If the statement were true, $\lim_{n \rightarrow \infty} 2^x / 2^y = 1$

$$\text{i.e. } \lim_{n \rightarrow \infty} 2^{x-y} = 1$$

$$\text{i.e. } \lim_{n \rightarrow \infty} x - y = 0.$$

But $x - y = \frac{\log n}{n} - \log \log n$. As $n \rightarrow \infty$, $\frac{\log n}{n} \rightarrow 0$ and $\log \log n \rightarrow \infty$. So $\lim_{n \rightarrow \infty} x - y \neq 0$

and hence the statement is false.

(e) For the given statement to be true, $\exists c \ni n^{10(\log \log n)^{100}} \leq c(\log n)!$ for all $n > N_0$ for some $N_0 \in \mathbb{N}$.

Also since $(\log n)! < (\log n)^{\log n}$ so for the given statement to be true, $\exists c \ni n^{10(\log \log n)^{100}} \leq c(\log n)^{\log n}$ for all $n > N_0$ for some $N_0 \in \mathbb{N}$

$$\text{But } \lim_{n \rightarrow \infty} \frac{n^{10(\log \log n)^{100}}}{(\log n)^{\log n}} = \infty$$

So, no such c exists. The given statement is false.

3. The Lucas Sequence 1, 3, 4, 7, 11, 18, 29, . . . is defined by $a_1 = 1, a_2 = 3, a_n = a_{n-1} + a_{n-2}$.

Prove that $a_n = O(1.75^n)$.

Sol: Given $a_n = a_{n-1} + a_{n-2}, \forall n \geq 3, a_1 = 1, a_2 = 3$.

Lets use generating functions to solve it.

Define, $A(x) = \sum_{i=1}^{\infty} a_i x^i \dots \dots (i)$

$$= x + 3x^2 + \sum_{i=3}^{\infty} a_i x^i$$

$$= x + 3x^2 + \sum_{i=1}^{\infty} (a_{i-1} + a_{i-2}) x^i$$

$$= x + 2x^2 + x \sum_{i=1}^{\infty} a_i x^i + x^2 \sum_{i=1}^{\infty} a_i x^i$$

$$= x + 2x^2 + xA(x) + x^2A(x)$$

Therefore, $A(x) = -\frac{x+2x^2}{1-x-x^2} \dots \dots (ii)$

The roots of the denominator are $\alpha = -\frac{(1+\sqrt{5})}{2}$ and $\beta = \frac{(-1+\sqrt{5})}{2}$

We can write $A(x)$ as $A(x) = -x(\frac{\gamma}{x-\alpha} + \frac{\delta}{x-\beta}) \dots \dots (iii)$

From (ii) & (iii) we get $\gamma = 1, \delta = 1$,

$$\begin{aligned}
\therefore A(x) &= -x\left(\frac{1}{x-\alpha} + \frac{\delta}{x-\beta}\right) \\
&= -x\left\{\frac{1}{\alpha(-\beta x-1)} + \frac{1}{\beta(-\alpha x-1)}\right\} \quad [\because \alpha\beta = 1] \\
&= x\left\{\frac{1}{\alpha(\beta x+1)} + \frac{1}{\beta(\alpha x+1)}\right\} \\
&= x\left\{\frac{1-x\beta+x^2\beta^2-x^3\beta^3+\dots}{\alpha} + \frac{1-x\beta+x^2\beta^2-x^3\beta^3+\dots}{\beta}\right\} \dots\dots (iv), \\
&\qquad\qquad\qquad \text{using binomial expansion of } (1+x)^{-1}
\end{aligned}$$

$$\begin{aligned}
\text{Equating (i) \& (iv), we get } a_n &= \frac{(-\beta)^n}{\alpha} + \frac{(-\alpha)^n}{\beta} \\
&= \frac{(1-\sqrt{5})^n + (1+\sqrt{5})^n}{2^n} \\
&< \frac{(-1.236)^n + (3.236)^n}{2^n} \\
&< (0.618)^n + (1.618)^n \\
&< 2(1.618)^n \\
&< 2(1.75)^n
\end{aligned}$$

Hence, $a_n = O((1.75)^n)$

Alternatively,

Given $a_1 = 1, a_2 = 3, a_n = a_{n-1} + a_{n-2}$, we shall show, $a_n \leq c(1.75)^n$

We proceed by induction on n .

Base Case : $a_1 \leq c, a_2 \leq 2c(1.75)^2$ valid for $c \geq 1$.

Inductive hypothesis : We assume, $a_i \leq c(1.75)^i$ For $i \leq c(1.75)^i$ for $i \leq n - 1$

$$\begin{aligned}
a_n &= a_{n-1} + a_{n-2} \\
&\leq c(1.75)^{n-1} + c(1.75)^{n-2} \\
&= c(1.75)^{n-2}(1 + 1.75) \\
&= c(1.75)^{n-2}(2.75) \\
&< c(1.75)^n \text{ for } c > 1.
\end{aligned}$$

Hence, $a_n = O((1.75)^n)$

4. For natural number p and q , the Ramsey number $R(p, q)$ is dened as the smallest integer n so that among any n people, there exist p of them who know each other, or there exist q of them who dont know each other. Note that $R(p, 1) = R(1, q) = 1$. Prove that:

- (a) $R(p + 1, q + 1) \leq R(p, q + 1) + R(p + 1, q)$
- (b) $R(p, q) \leq \binom{p+q-2}{p-1}$

Sol: Let us observe the problem in terms of graph theory. We denote $G(V; E)$ an undirected graph as follows,

$$V = \{v \mid v \text{ is a person}\}$$

$$E = \{(u, v) \mid u \text{ knows } v\}$$

Then by definition, $R(p, q)$ is the smallest natural number n such that a graph with n vertices has either a size- p clique or a size- q independent set.

(a) Let
$$R(p+1, q) + R(p, q+1) = n$$

We need to prove that for any graph G on n vertices, G either has a size- $(p + 1)$ clique or a size- $(q + 1)$ independent set. Let us choose any vertex $v \in V$. We define,

$$N_v = \{u \in V \mid (u, v) \in E\}$$

$$\overline{N}_v = \{v \in V \mid (u, v) \notin E\}$$

Now either $|N_v| \geq R(p, q + 1)$ or $|\overline{N}_v| \geq R(p + 1, q)$. Because if not then

$$\begin{aligned} |V| &= |N_v| + |\overline{N}_v| + |\{v\}| \\ &\leq R(p, q + 1) - 1 + R(p + 1, q) - 1 + 1. \\ &< n \end{aligned}$$

a contradiction.

CASE I

If $|N_v| \geq R(p, q+1)$ then the induced subgraph of G on the vertex set N_v has a size- p clique or a size- $(q+1)$ independent set. Thus this subgraph along with vertex v has either a size- $(p+1)$ - clique or a size- $(q+1)$ independent set. Hence G will have either a size- $(p+1)$ clique or a size- $(q+1)$ independent set.

CASE II

If $|\overline{N}_v| \geq R(p+1, q)$ then the induced subgraph of G on the vertex set \overline{N}_v has a size- $(p+1)$ clique or a size- q independent set. This subgraph along with vertex v has either a size- $(p+1)$ clique or a size- $(q+1)$ independent set.

Hence G will have either a size- $(p+1)$ clique or a size- $(q+1)$ independent set.

Hence we have proved by exclusive case studies that

$$R(p + 1, q + 1) \leq R(p, q + 1) + R(p + 1, q)$$

(b) We prove the given inequality by induction on p and q . Let

$$R(i, q) \leq \binom{i+q-2}{i-1}$$

$$R(p, i) \leq \binom{p+i-2}{p-1}$$

Base Case: $R(1,q)=1 \forall p \in \mathbb{N}$

$R(p,1)=1 \forall q \in \mathbb{N}$

Lets assume that the given inequality holds for $p=k$ and $q=l$, and prove it for $p=k+1$ and $q=l+1$.

$$\begin{aligned} R(k+1,l+1) &\leq R(k,l+1) + R(k+1,l) \quad \text{from 4(a)} \\ &\leq \binom{k+l-1}{k-1} + \binom{k+l-1}{k} \\ &= \binom{k+l}{k} \quad [\because \binom{n}{k} + \binom{n}{k-1} = \binom{n+1}{k}] \end{aligned}$$

Hence, by induction, the given inequality is true.

5. Write the negation of the following statement: $\forall x \leq 0 \exists y \in \mathbb{N} (y \geq x) \wedge (y \text{ is a prime})$

$$\begin{aligned} \text{Sol: } &\neg (\forall x \leq 0 \exists y \in \mathbb{N} (y \geq x) \wedge (y \text{ is a prime})) \\ &= \exists x \geq 0, \forall y \in \mathbb{N}, \neg \{(y \geq x) \wedge (y \text{ is a prime})\} \\ &= \exists x \geq 0, \forall y \in \mathbb{N}, (y \leq x) \vee (y \text{ is a prime}) \end{aligned}$$

6.

(a) We know that $\sqrt{3}$ is not rational. Using this prove that $\sqrt{3} + \sqrt{24}$ is not rational.

(b) If m is a positive integers such that \sqrt{m} is not rational, then prove that for any positive integer n the number $\sqrt{m} + \sqrt{n}$ is not rational.

Sol:

(a) Let's assume that $\sqrt{3} + \sqrt{24}$ is rational. Then,

$$\sqrt{3} + \sqrt{24} = \frac{p}{q}, p, q \in \mathbb{Z}$$

or

$$q\sqrt{24} = p - \sqrt{3}q$$

or

$$24q^2 = (p - \sqrt{3}q)^2$$

or

$$24q^2 = p^2 + 3q^2 - 2\sqrt{3}pq$$

or

$$\sqrt{3} = \frac{p^2 - 21q^2}{2pq} \quad \text{which is rational.}$$

But $\sqrt{3}$ is not rational. Thus we get a contradiction.

(b) Given, \sqrt{m} is not rational, we assume, if possible, that $\sqrt{m} + \sqrt{n}$ is rational. Then

$$\sqrt{m} + \sqrt{n} = \frac{p}{q} \quad \text{where } p, q \in \mathbb{Z}$$

or

$$\sqrt{n}q = p - \sqrt{m}q$$

or

$$m^2q^2 - 2\sqrt{m}pq + p^2 = nq^2$$

or

$$\sqrt{m} = \frac{p^2 - 21q^2}{2pq} \quad \text{which shows that } \sqrt{m} \text{ is rational, a contradiction.}$$

Alternatively,

Assume for the sake of contradiction that $\sqrt{m} + \sqrt{n}$ is rational

$$(\sqrt{m} + \sqrt{n})(\sqrt{m} - \sqrt{n}) = m - n.$$

This shows that $(\sqrt{m} - \sqrt{n})$ is rational, from the fact that a rational no. multiplied by an irrational no. cannot give a rational no. Now, if $(\sqrt{m} + \sqrt{n})$ and $(\sqrt{m} - \sqrt{n})$ are both rationals, their sum, $2\sqrt{m}$ must also be rational, a contradiction.