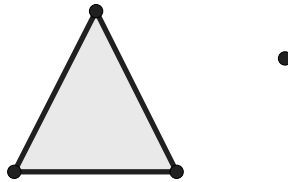


Quiz-2

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1. If G is a graph with even number of vertices and each vertex has even degree (that is, even number of neighbors) then G has even number of edges. Either prove the statement or disprove it by demonstrating a counter example.

Solution:



Number of vertices of the graph is 4 but number of edges is 3 which is odd. Hence we get a contradiction.

2. Is it possible to draw 9 line segments in the plane \mathbb{R}^2 such that each line segment intersects exactly 3 other line segments.

Solution:

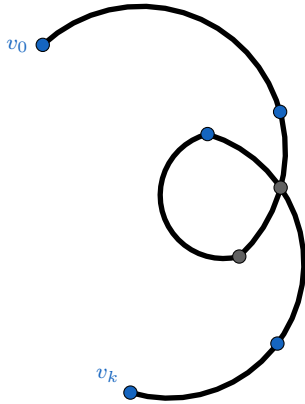
Consider each line segment as a vertex of a graph G and draw an edge between each vertex whose corresponding line segment intersects each other.

Then $\deg(v) = 3 \forall v \in G$. Then $|E| = \frac{3 \times 9}{2} \notin \mathbb{Z}$ where $|E|$ denotes number of edges. Hence the above statement cannot be possible.

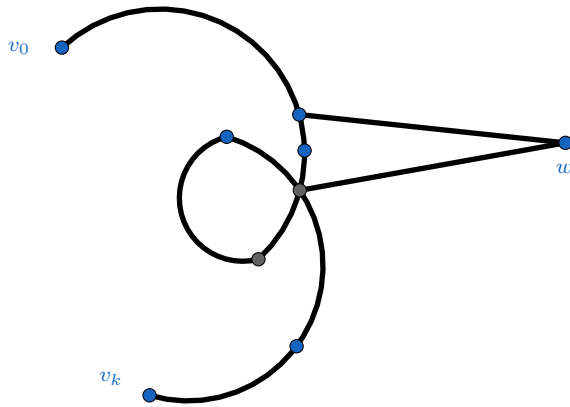
3. Prove that a graph has an Eulerian path if and only if the number of vertices in the graph with odd degree is 0 or 2.

Solution:

First part is fairly straightforward.



Let v_0, \dots, v_k be an Euler trail. So each v_i has even degree except v_0 and v_k . If this is an Euler trail then v_0 and v_k are two endpoints; hence number of odd degree vertices are two. If $v_0 = v_k$ then it is an Euler tour. There the number of odd degree vertices are zero.



Conversely, let l be the longest path in G . There are at most two odd degree vertices. Let v_0, \dots, v_k be the path. v_k is not adjacent to any other vertex because if it was the case then l will not be the longest path.

Let l be not Eulerian then there exist some path v_i, v_{i+1} which are not part of l because each v_i is of even degree. So there exist another path which is also not part of l . If G has zero odd degree vertex then $v_k = v_0$. Now if we start a path from v_i then traverse l that will be longer than l . (Contradiction)

So, l is Eulerian.

Now let G has 2 odd degree vertices say u and v . Add new vertex w . Join uw and vw . So degree of w is 2. So degree of v and u is even.

So $H = G \cup w$. So H has all even vertices of even degree. So H has an Euler tour. Choose an Euler tour that start and end in w . If we remove w from the Euler tour we get Euler trail. Hence proved.

4. Consider an $n \times m$ rectangular grid, with the co-ordinates of the corners being

$(0, 0)$, $(n, 0)$, (n, m) and $(0, m)$. How many paths are there along the rectangular grid from $(0, 0)$ to (n, m) such that

- (a) The paths are shortest among all paths from $(0, 0)$ to (n, m) .
- (b) Every horizontal move of unit length is followed by at least one vertical move and the paths are the shortest.

Solution:

- (a) The number of shortest paths is $\binom{n+m}{n}$.

6. For a simple graph G an independent set is a subset S of vertices such that no two vertices in the set S are adjacent. Prove that

- (a) If a graph on n vertices has a matching of size k then G cannot have an independent set of size $> (n - k)$.
- (b) If the graph G on n vertices has a maximal matching of size k then G has an independent set of size $\geq (n - 2k)$.

Solution:

- (a) Let there be a matching of size k which is forming a bipartite graph component X and Y . So there are $n - 2k$ vertices remaining in G . Let S be the set consisting of those remaining vertices of G . This set is independent.

Let S is only adjacent to Y . So $S + X = S'$ is also independent.

$$|S'| = n - 2k + k = n - k$$

But we cannot add any vertex to S because all other vertices are adjacent to S . Hence, G cannot have independent set of size $> n - k$.

- (b) Let S be a maximal matching of size k . S' be a set consists of remaining $n - 2k$ vertices. If this set is not independent, then there exists some edge between vertices v_1, v_2 which belongs to S' , so then we extend the size of S . But S is maximal hence a contradiction.