

## Lecture 17: Graph Theory

Instructor: Sourav Chakraborty

Scribe: Shashank S. Bhavishaya

## 1 Graph Theory

**Definition 1** A **Graph** is represented by  $G(\mathbf{V}, \mathbf{E})$ , whereas  $\mathbf{V}$  stands for set of vertices of graph and  $\mathbf{E}$  stands for set of edges ( $\mathbf{E} \subseteq \mathbf{V} \times \mathbf{V}$ ).

- A **Simple Graph** has no multiple edges, no self loop and is undirected.
- $\forall v \in V, N_{gb}(v) = \{w : (v, w) \in E\}$ .
- $|N_{gb}(v)| = \deg(v) = d(v)$ .

For undirected graph  $(v, w)$  and  $(w, v)$  are same.

- Degree sequence is sequence of degree of vertices, labelling the vertices gives permutation of degree sequences.
- So without loss of generality **degree sequence** of a graph is  $(d_1, d_2, d_3, \dots, d_n)$  whereas  $(d_1 \leq d_2 \leq d_3 \leq \dots \leq d_n)$

**Theorem 1.1** For every graph  $G(V, E)$

$$\sum d(v) = 2|E|$$

**Corollary 1.2** If  $(d_1, d_2, \dots, d_n)$  is a degree sequence of a graph then,

$$\sum_{i=1}^n d_i = \text{even}$$

**Corollary 1.3** Number of odd degree vertices in graph is always even.

$$\sum_{i=1}^n d_i = \sum_{\text{even}} d_e(v) + \sum_{\text{odd}} d_o(v) = \text{even}.$$

- $(d_1, d_2, \dots, d_n)$  is a valid degree sequence of a simple undirected graph, then each of  $d_i$ 's will satisfy :  $0 \leq d_i \leq n-1$ .

- Let we have ordered degree sequence such that first (i-1) vertices has degree 0 and  $d_i \neq 0$ .

**Theorem 1.4** *In every connected graph there exist atleast two vertices with same non-zero degree.*

*Proof.*

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Let H be connected subgraph of  $G(V, E)$  with n vertices.

$\therefore$  vertices in H is n-(i-1) and every degree is  $1 \leq d_i \leq (n - (i - 1)) - 1 = n - i$

$\therefore$  total number of vertices in H is n-(i-1) and degree is between 1 to n-i,

$\therefore$  by pigeon hole principle(PHP), at least 2 has same degree. □

**Theorem 1.5**  $\Delta G \leq$  total number of vertices with non zero degree.

*Proof.* by pigeon hole principle(PHP) as described above. □

( $\Delta G$  is defined as maximum degree of a vertex in graph).

## 2 Graph Isomorphism

Given two different graph  $G_1$  and  $G_2$  such that,

$G_1(V_1, E_1)$

$G_2(V_2, E_2)$

these two graph are **isomorphic** if their is a function 'p' on  $V_1$  and  $V_2$

$p : V_1 \rightarrow V_2$  (bijection)

such that,  $(v, w) \in E_1 \Leftrightarrow (p(v), p(w)) \in E_2$

If degree sequence of  $G_1 \neq$  degree sequence of  $G_2$  then  $G_1$  is **not isomorphic** to  $G_2$ .

## 3 Tree

**Definition 2** *A maximum connected acyclic graph is tree.*

*Number of edges is n-1 with total number of vertices is n.*

*The following are equivalent*

- *A tree is minimally edge connected graph.(removing any single edge will disconnect graph).*

- A tree is maximally acyclic graph.(including any single edge will create a cycle).
- A tree is a connected graph with  $n-1$  edges.
- $\forall u,v,\exists$  a unique path from  $u$  to  $v$

**Lemma 3.1** If  $G$  is a tree then there exist two vertices with degree 1. if not then there exist a cycle and it is not a tree.

## 4 Connected Component

If  $G = H_1 \cup H_2 \cup \dots \cup H_n$  whereas,  
 $H_1(V_1, E_1), H_2(V_2, E_2), \dots, H_n(V_n, E_n)$  such that,  
 $V_1, V_2, \dots, V_n$  is partition of  $V$ .  
 $E_1, E_2, \dots, E_n$  is partition of  $E$ .

- $H_i$  is connected.
- $\forall (u,v), u \in V_i, v \in V_j (j \neq i)$  then  $(u,v) \notin E$ .

### 4.1 k-connected

Graph is **2-connected** if  $\forall v, G \setminus v$  is connected.

Graph is **2-edge-connected** if  $\forall e, G \setminus e$  is connected.

Graph is **k-connected** if  $\forall (v_1, v_2, \dots, v_{k-1}), G \setminus (v_1, v_2, \dots, v_{k-1})$  is connected.

## 5 Bipartite graph

**Definition 3** A graph  $G(V,E)$  in which  $V$  can be partitioned into two subsets such that there is no edge between vertices of any one set.

**Property :**

- There exists no cycle with odd number of vertices.
- Partition  $V = V_1 \cup V_2, \text{s.t. } E \subseteq V_1 \times V_2$