

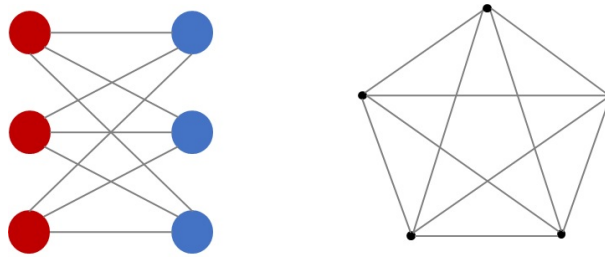
Lecture 21: Planar Graph

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1 Planarity

1.1 Observations

Observation 1.1 K_5 or $K_{3,3}$ cannot be drawn on a plane such that there is no intersection of two edges.

Figure 1: $K_{3,3}$ and K_5

Lemma 1.2 If G can be drawn (no two edges intersect) on $\mathbb{R}^2 \implies G$ can be drawn on a sphere.

Proof :

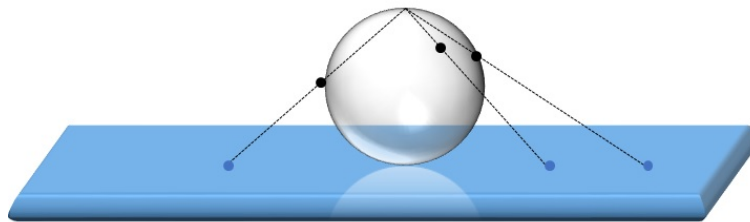


Figure 2: every point on the sphere has a bijection to a point on the plane

As every point on the sphere has bi-jection to a point on the plane.

- The number of points on the sphere is one more than the number of points on the plane.

1.2 Operations

- If a graph is planar, then if some vertices or edges be removed then the resulting graph would also be planar.
- If a graph is planar, then if some adjacent vertices be collapsed then the resulting graph would also be planar.

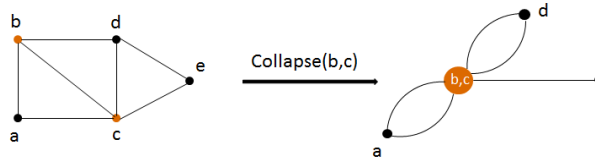


Figure 3: A planar graph after collapsing two vertices remains planar

How to check if a graph is planar?

1. Delete a vertex
2. Delete an edge
3. Contract an edge(Collapse two adjacent vertices).

if \exists a sequence of moves 1,2,3 such that $G \rightarrow H$ then we say G has a minor H .

If a graph G has K_5 or $K_{3,3}$ as minor then G is not planar.

1.3 Theorems

Theorem 1.3 A graph G is planar $\iff G$ has no K_5 or $K_{3,3}$ as minor. (Kuratowski's theorem).

Proof. Next lecture □

Theorem 1.4 If G is a connected planar graph then $|V| + |E| - |F| = 2$, where $|V|$: number of vertices, $|E|$: number of edges, $|F|$: number of faces of the graph.

Proof.

Induction on number of edges of the graph, i.e. $|E|$.

Base Case : When $|E| = 1$, as the graph is connected, so when there is a single edge means there are only 2 vertices.

$$|V| = 2, |E| = 1, |F| = 1.$$

so, $|V| + |E| - |F| = 2 + 1 - 1 = 2$ (proved).

Induction Hypothesis : For a planar connected graph with number of edges $\leq k$ the given statement is true.

Inductive step : Let G be a planar connected graph with number of edges $= k+1$. Every edge in a cycle is a part of 2 faces.

\exists a cycle $\implies \exists$ edge e , such that e is a part of 2 faces.

Let G' be a sub-graph of G such that $G' = G - e$.

G' is connected as e is a part of a cycle and G' is also planar.

Then, $(|V(G')| = |V|, |E(G')| = |E| - 1, |F(G')| = |F| - 1)$.

According to I.H. $|V(G')| - |E(G')| + |F(G')| = 2$

$$|V| - (|E| - 1) + (|F| - 1) = 2$$

$$|V| + |E| - |F| = 2 \text{ (Proved).}$$

□

1.4 Definitions

Definition 1.5 *Maximal planar graph : It is a planar graph and if any more edge is introduced then it will not be planar anymore.*

1.5 Lemmas

Lemma 1.6 *For a planar graph, $|E| \leq 3|V| - 6$.*

Proof : Every edge in a maximal planar graph is a part of a cycle and as a result part of 2 faces. So $\sum_i |E(f_i)| = 2|E|$.

Now as every face contains at least 3 edges, so $3|F| \leq \sum_i |E(f_i)|$.

Equating the above two equations, $3|F| \leq 2|E|$.

$$\text{so, } |F| \leq \frac{2}{3}|E|.$$

From Euler's formula, $2 = |V| - |E| + |F|$

$$2 \leq |V| - |E| + \frac{2}{3}|E|$$

$$2 \leq |V| - \frac{1}{3}|E|$$

$$|V| - 2 \geq \frac{|E|}{3}$$

$$|E| \leq 3|V| - 6$$

Lemma 1.7 *For a bipartite planar graph, $|E| \leq 2|V| - 4$.*

Proof : As there is no odd cycle in a bipartite graph so every edge in a bipartite planar graph is a part of a 4-cycle and part of 2 faces. So $\sum_i |E(f_i)| = 2|E|$.

Now as every face contains at least 4 edges, so $4|F| \leq \sum_i |E(f_i)|$.

Equating the above two equations, $4|F| \leq 2|E|$.

$$\text{so, } |F| \leq \frac{1}{2}|E|$$

From Euler's formula, $2 = |V| - |E| + |F|$

$$2 \leq |V| - |E| + \frac{1}{2}|E|$$

$$2 \leq |V| - \frac{1}{2}|E|$$

$$|V| - 2 \geq \frac{|E|}{2}$$

$$|E| \leq 2|V| - 4$$