

Lecture 22: Proof Of Kuratowski's theorem

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1 Preliminaries

1.1 Definitions

1. **Planar Graph** : A planar graph is a graph that can be embedded in the Euclidean plane, i.e., it can be drawn on the plane in such a way that no edges cross each other. Such a drawing is called a **plane graph** .
2. K_5 is a non-planar graph with minimum number of edges on 5 vertices and $K_{3,3}$ is a non-planar graph with minimum number of edges on 6 vertices.
3. A **minimal non-planar** graph is a non-planar graph G such that every proper sub-graph of G is planar.
4. **Minor of a Graph** : An undirected graph H is a minor of another undirected graph G if H can be obtained from G by
 1. contracting some edges
 2. Deleting some edges
 3. Deleting some vertices.

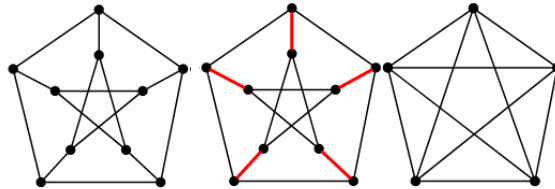


Figure 1: Petersen graph has K_5 as minor. Red edges are contracted to get K_5 .

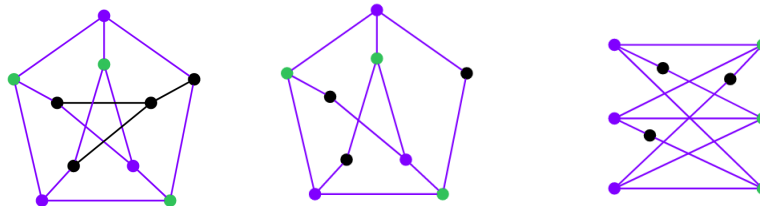


Figure 2: Petersen graph has $K_{3,3}$ as minor. Black edges are deleted to get $K_{3,3}$.

1.2 Lemmas

1. **Lemma 1** For every face of a given plane graph G , there is a drawing of G for which the face is exterior.

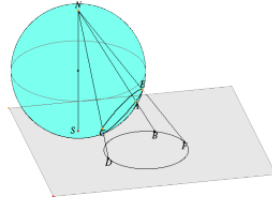


Figure 3: We draw the graph on a sphere, and then project it from a point on the face f

2. **Lemma 2** Every minimal non-planar graph is 2-connected.

proof Assume, for contradiction, that there exists a minimal non-planar graph $G = (V, E)$ that is not 2-connected.

There exists a vertex v whose removal disconnects G . Let C_1, C_2 be the components of G/v . By the minimality of G , the induced subgraph on $C_1 \cup \{v\}$ and $C_2 \cup \{v\}$ are both planar. We can embed both graphs with v on the unbounded face, and merge both copies of v to get a planar graph. This is contradiction.

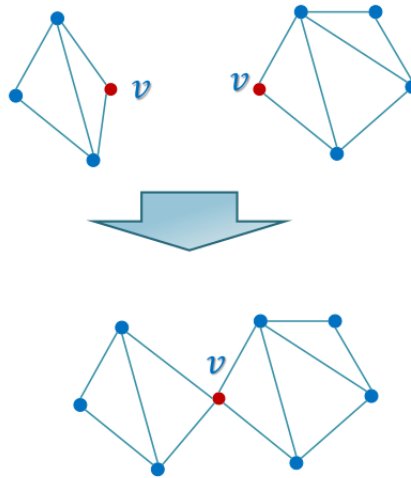


Figure 4: resulting graph is again planar

□

3. **Lemma 3** Let $G = (V, E)$ be a graph with fewest edges among all non-planar graphs

without a K_5 or $K_{3,3}$ as minor. Then G is 3-connected.

proof Assume, for contradiction, that there exists a minimal non-planar graph $G = (V, E)$ without a K_5 or $K_{3,3}$ as minor is not 3-connected. There exists a pair of vertices u, v whose removal disconnects G . Let C_1, C_2 be the components of $G - \{u, v\}$. By the minimality of G , the induced subgraph on $C_1 \cup \{u, v\}$ and $C_2 \cup \{u, v\}$ are both planar. Now there arises two cases:

- (a) **Case 1** : Edge (u, v) is present in the graph G . Then $C_1 \cup \{u, v\}$ and $C_2 \cup \{u, v\}$ both are planar, by minimality of G . Now we can bring the edge $\{u, v\}$ on the outer face of $C_1 \cup \{u, v\}$ and $C_2 \cup \{u, v\}$ using lemma 1. Since both graphs are planar, by joining vertices u in $C_1 \cup \{u, v\}$ to u in $C_2 \cup \{u, v\}$ and vertices v in $C_1 \cup \{u, v\}$ to v in $C_2 \cup \{u, v\}$ creates a planar graph. Contradiction to our assumption that G is non planar.

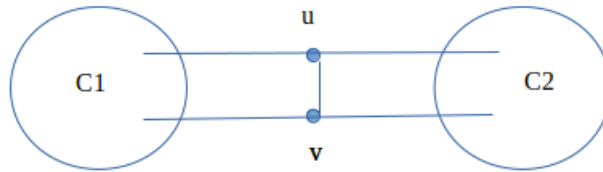


Figure 5: Graph G which is minimal non-planar graph

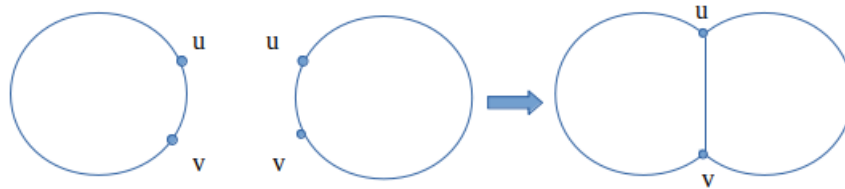


Figure 6: resulting graph is again planar

- (b) **Case 2** : Edge $\{u, v\}$ is not present. Then $C_1 \cup \{u, v\}$ and $C_2 \cup \{u, v\}$ both are planar, by minimality of G . Since u, v are not connected by direct edge, by bringing the vertex u on the outer face of $C_1 \cup \{u, v\}$ might not bring the v on outer face of $C_1 \cup \{u, v\}$ simultaneously and vice versa. So the trick we apply is to join u and v by an edge temporarily and then we can bring both the vertices u, v on the outer face.

Lets call $G_1 = C_1 \cup \{u, v\} \cup (u, v)$ and $G_2 = C_2 \cup \{u, v\} \cup (u, v)$. Here we cannot guarantee that G_1 and G_2 are planar. Because we have added an extra edge (u, v) . But are we doing any mistake by adding extra edge? Does this

introduce K_5 or $K_{3,3}$ in the graph $C1 \cup \{u,v\}$ or $C2 \cup \{u,v\}$? The answer is No. Because if it would have introduced a K_5 or $K_{3,3}$ in the graph $C1 \cup \{u,v\}$ or $C2 \cup \{u,v\}$ that means the path joining u and v through graph $C2 \cup \{u,v\}$ contains a K_5 or $K_{3,3}$. This implies the original graph G contains a K_5 or $K_{3,3}$ as minor. Which is contradiction, because we started with a non planar graph G which doesn't contain K_5 or $K_{3,3}$ as minor.

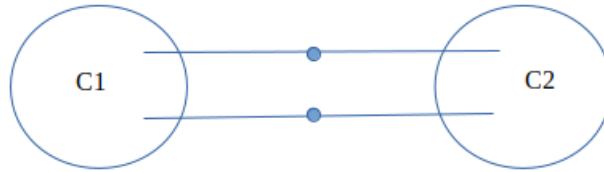


Figure 7: resulting graph is again planar

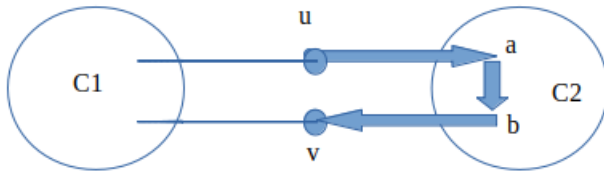


Figure 8: We can contract the path $u \rightarrow a \rightarrow b \rightarrow v$ to make an edge. And if $C1 \cup \{u,a,b,v\}$ contains K_5 or $K_{3,3}$ then G has K_5 or $K_{3,3}$.

□

1.3 Proof

Theorem 1.1 (Kuratowski's Theorem) *A graph is planar iff it does not have K_5 or $K_{3,3}$ as minors.*

proof We know that if a graph contains K_5 or $K_{3,3}$ as a minor graph, then it is not planar.

It remains to prove that every non-planar graph contains K_5 or $K_{3,3}$ as minor.

Proof Strategy: For proving this

1. It suffices to prove this only for minimal non- planar graphs.
2. We will show that every minimal non-planar graph with no K_5 or $K_{3,3}$ as minor must be 3-connected.

3. We then show that every 3-connected graph with no K_5 or $K_{3,3}$ as minor is planar. But we started with a non-planar graph. which is Contradiction! So a non-planar graph must contain K_5 or $K_{3,3}$ as minor graph.

We need to show that if a graph is non-planar then it must contain a K_5 or $K_{3,3}$ as minor graphs. Let G be the smallest graph in the set of non planar graph which is 3-connected. Then we prove the contra-positive of statement "G doesn't have K_5 and $K_{3,3}$ as minor and G is non-planar" i.e. "G is planar or G has K_5 or $K_{3,3}$ as minor".

1. **Case 1** : Let denote $N(v)$ as set of all vertices neighbour to v . Then if $|N(u) \cap N(v)| \geq 3$ then G has a K_5 a minor. Which is contradiction.

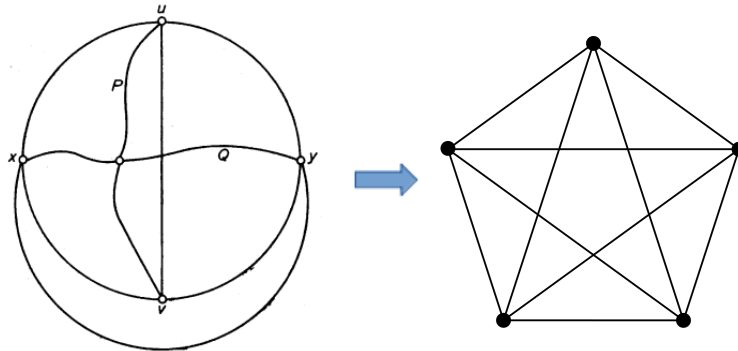


Figure 9: This graph has K_5 as minor.

2. **Case 2** : Here $|N(u) \cap N(v)| \leq 2$ is satisfied but it is also non planar. Let
 $N(u) = a_1, a_2, a_3, \dots, a_n$
 $N(v) = b_1, b_2, b_3, \dots, b_n$
 If the neighbour of u and v interleave then G has a $K_{3,3}$ minor. Which is contradiction.

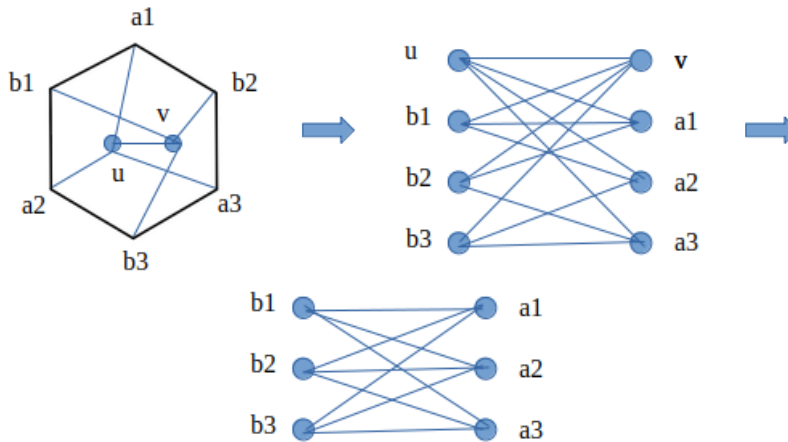


Figure 10: This graph has $K_{3,3}$ as minor.

3. **Case 3** : If the neighbour of u and v doesn't interleave and $|N(u) \cap N(v)| \leq 2$ is satisfied then G is planar.

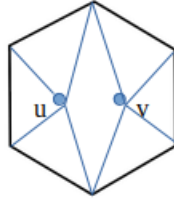


Figure 11: This is a planar graph

All the cases gives contradiction. This proves the theorem. □

1.4 References

1. <http://www.math.caltech.edu/2014-15/2term/ma006b/10Planar3.pdf>
2. https://en.wikipedia.org/wiki/Graph_minor
3. <https://www.math.hmc.edu/kindred/cuc-only/math104/lectures/lect17-slides-handout.pdf>
4. Graph theory with applications by J. A. Bondy and U. S. R. Murty